

BTAC  
与  
机场安全模型

Lenx Wei & Xiaoning Li

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# BTAC与机场安全模型

两个往返于机场的奶爸对内存安全的思考

Lenx Wei & Xiaoning Li



# 源起

内存漏洞伴随可寻址内存的产生而产生

- \* C/ASM/C++... 传统系统编程语言缺乏内存安全机制
  - \* buffer overflow/UAF/Random RW/...
- \* 人脑无法掌控Raw pointer引发的复杂性爆炸
  - \* C语言魔咒
- \* 产品Benchmark大战加剧了过度优化和指针滥用
  - \* Benchmark魔咒

# 洞是挖不完的 防御者的噩梦

- \* C语言魔咒/Benchmark魔咒
  - \* 别忘了编译器、ld、lib、JVM、Driver、Kernel等等
- \* Exploit Mitigation漏洞缓解技术日益获得重视
  - \* Stack Cookie
  - \* DEP
  - \* ASLR
  - \* EMET的努力
- \* 但是，还不够...
  - \* Pwn2Own每年打脸...
  - \* iOS Jailbreak...

# CFI的提出

- \* Control-Flow Integrity Principles, Implementations, and Applications, *Martin Abadi, Mihai Budiu, Ulfar Erlingsson, Jay Ligatti*, CCS'05
- \* 维持程序里原有的控制流完整性，防止控制流注入和任意劫持

# CFI的6年沉默

## \* Performance overheads

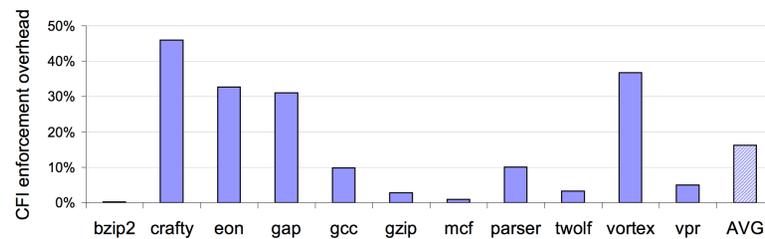
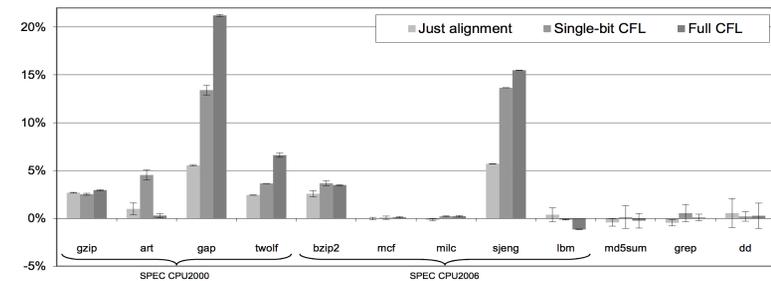


Figure 4: Execution overhead of inlined CFI enforcement on SPEC2000 benchmarks.



## \* Compatibility nightmare

### \* Legacy libraries

# CFI的再起

- \* FPGate by PKU&BitBlaze teams(2012)
  - \* Protect Function pointers, virtual methods and so on
  - \* Provide a binary compatibility solution for PE files
  - \* **Max: 1.41%, Avg: 0.36%**
  - \* MS Bluehat Prize Contest Special Recognition Award
- \* Followed by
  - \* Control-flow Guard, MS, 2014
  - \* Enforcing Forward-Edge Control-Flow Integrity in GCC & LLVM, Google, 2014

# CFI的复兴

- \* CCFIR, S&P'13
- \* CFI for COTS, Usenix Security'13
  
- \* KCoFI, S&P'14
- \* MCFI, PLDI'14
  
- \* OCFI, NDSS'15
- \* CCFI, CCS'15

# CFI的困境

- \* OoC, S&P'14
- \* DVE, by Yuange
- \* CFB, Control-flow bending, Usenix Security'15
- \* Control Jujutsu, CCS'15

# To be CFI, or not to be?

- \* 细粒度不是CFI的未来
  - \* Performance overhead
  - \* Compatibility issues
  - \* 内存里的猪队友
    - \* DVE attacks
- \* 何去何从?

# CFI角色的反思

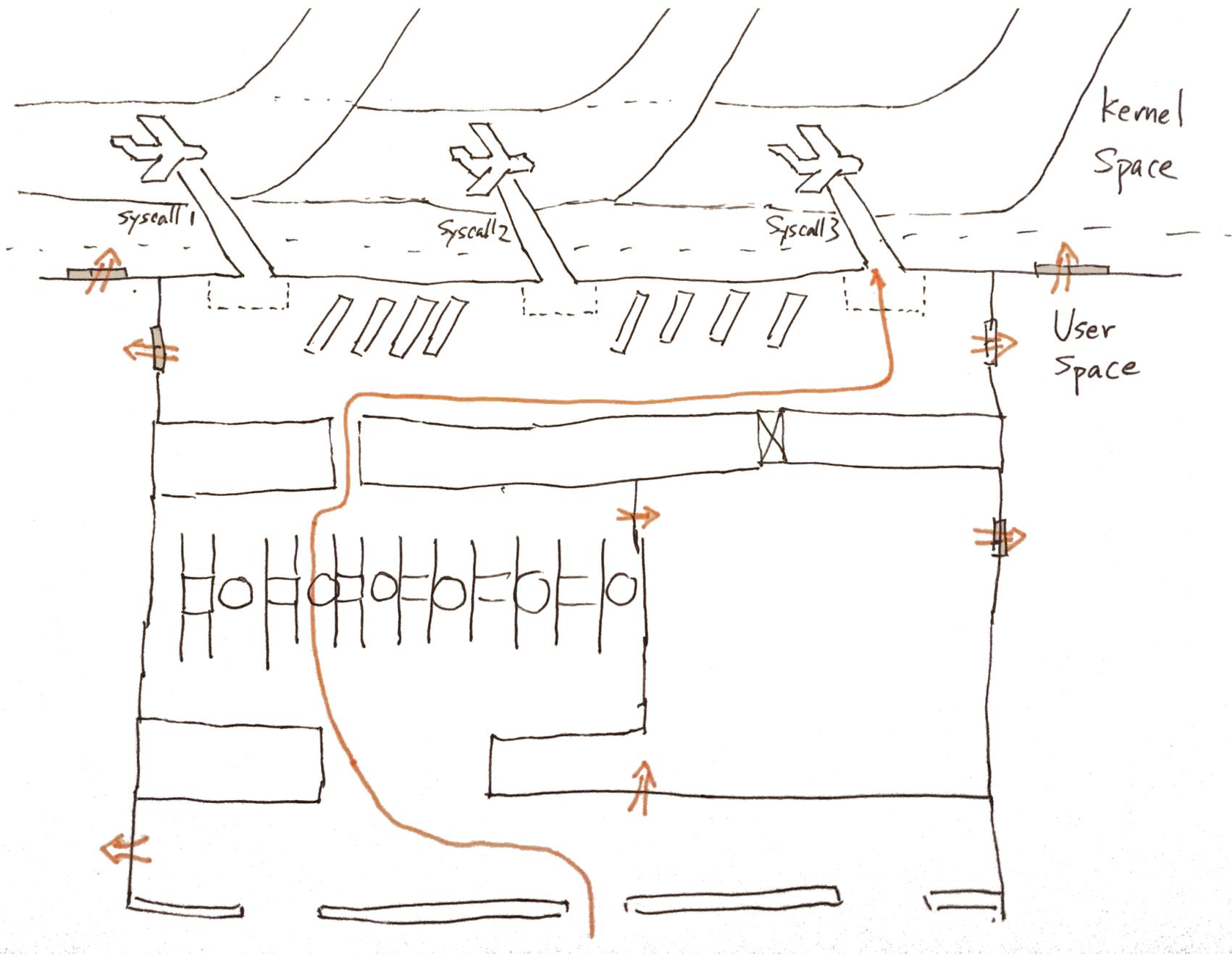
- \* Con:

- \* CFI不可能完成所有的防护

- \* Pro:

- \* CFI增加了一个新的防护层，一个新的基石

# 机场安全类比



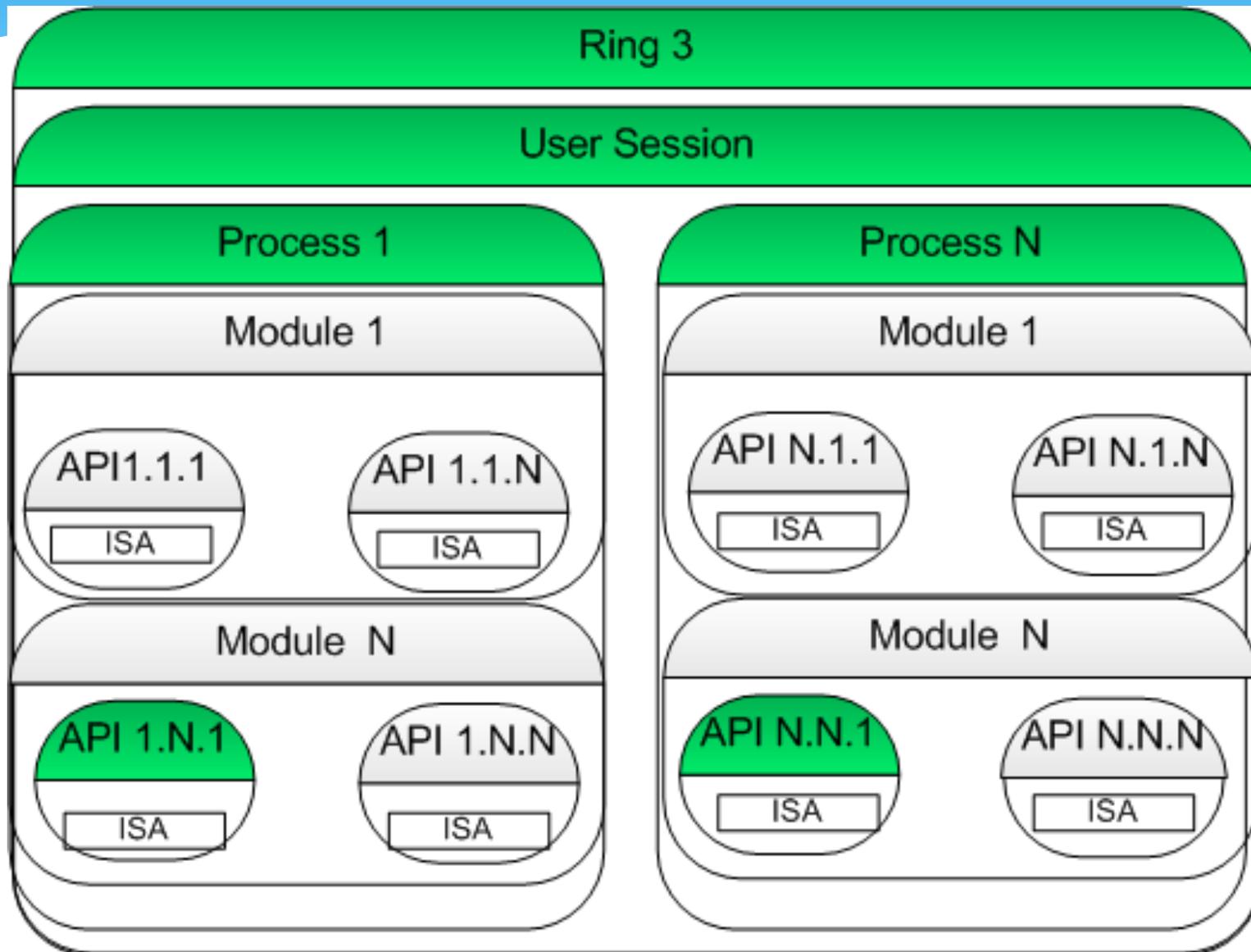
# 机场安全模型

- \* CFI保证了显性的执行通道
  - \* 攻击者很难击穿墙壁，绕过安全门
  - \* 攻击者依然可以选择所有允许的通道
- \* 关键点(API/Syscall)必须要可靠的Access Control
  - \* 每个访问的授权凭证(Passport/Boarding card)
  - \* 关键检查点的不可绕过可以依靠CFI来保证

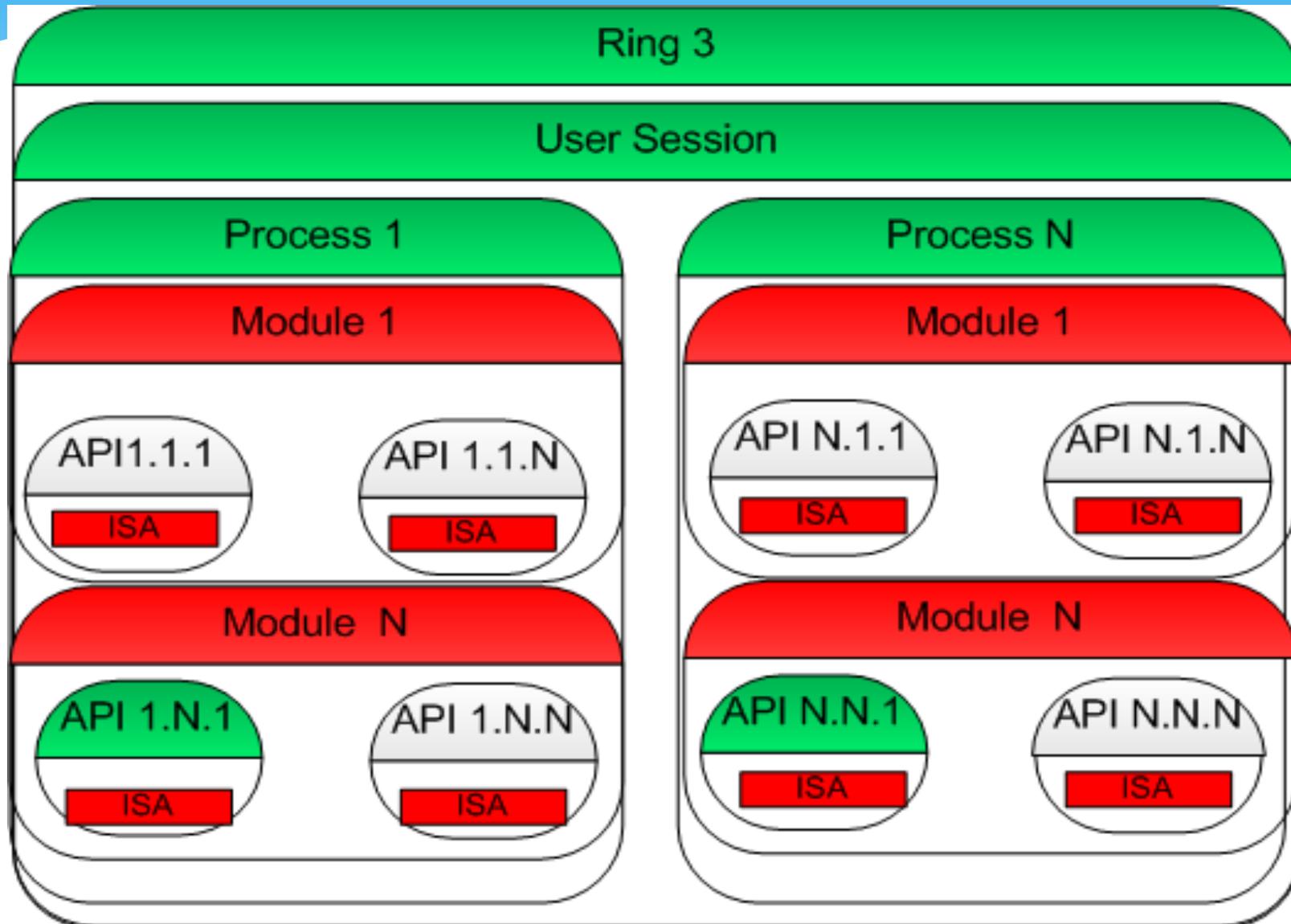
# Existing System Access Controls

- \* User level
  - \* Resource access per user permission
- \* Process Level
  - \* Resource access per process integrity level
- \* API Level
  - \* API resource access per HIPS policy

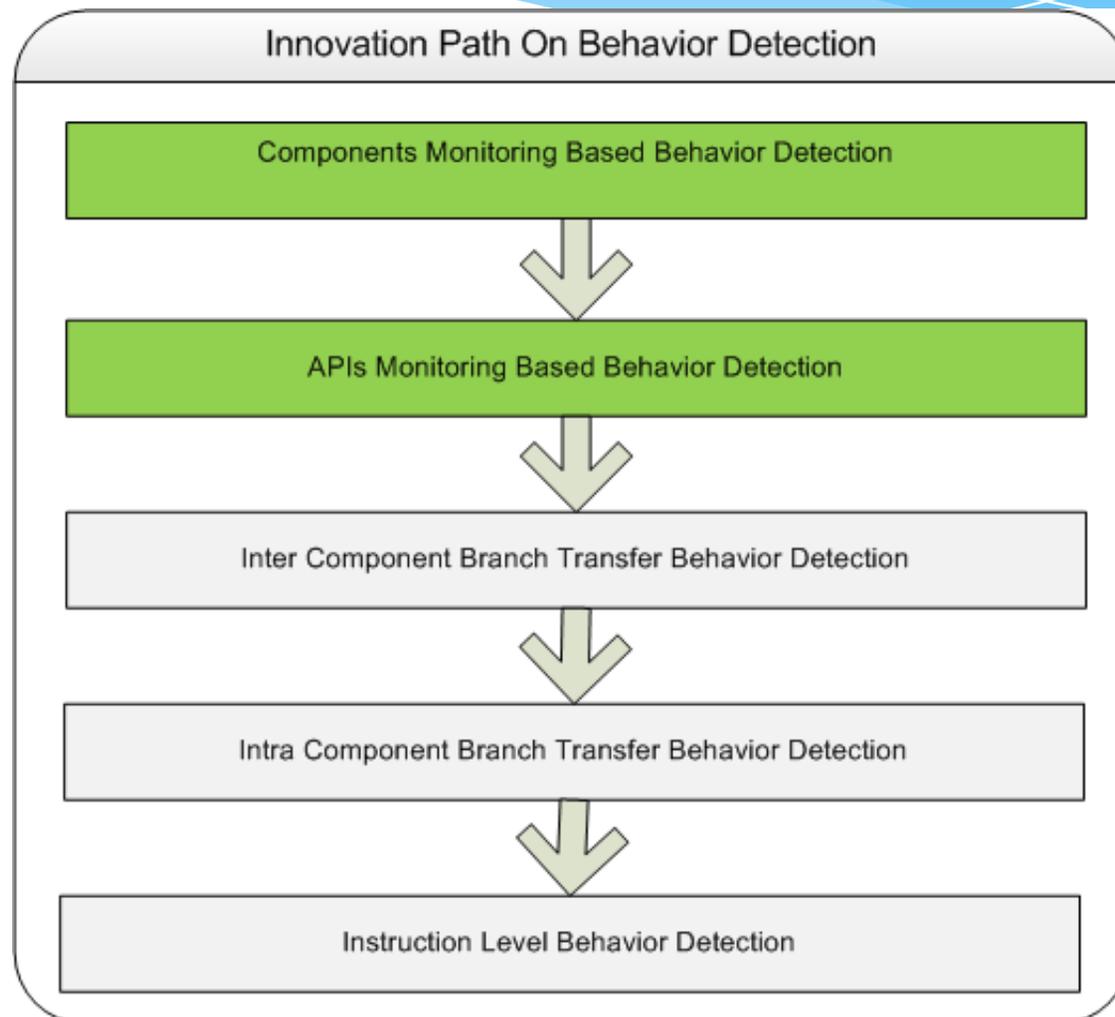
# Granularity from OS Ring 3 view



# Module/ISA Level granularity gap

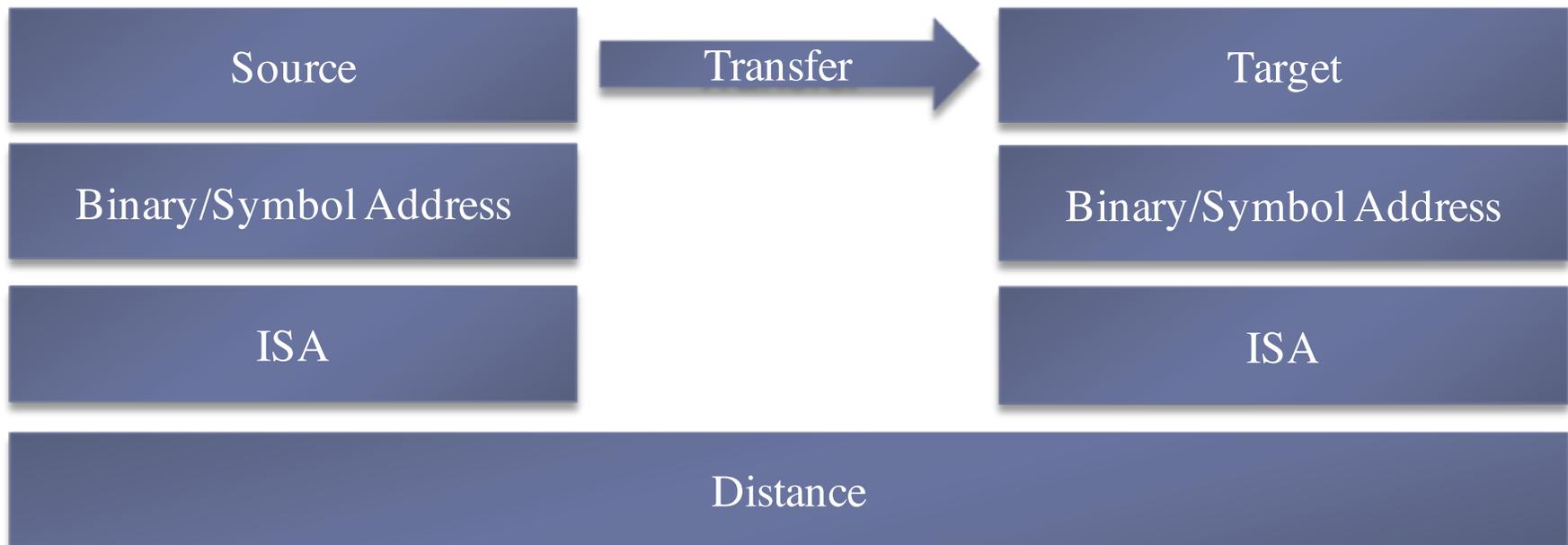


# 需要更细的AC粒度



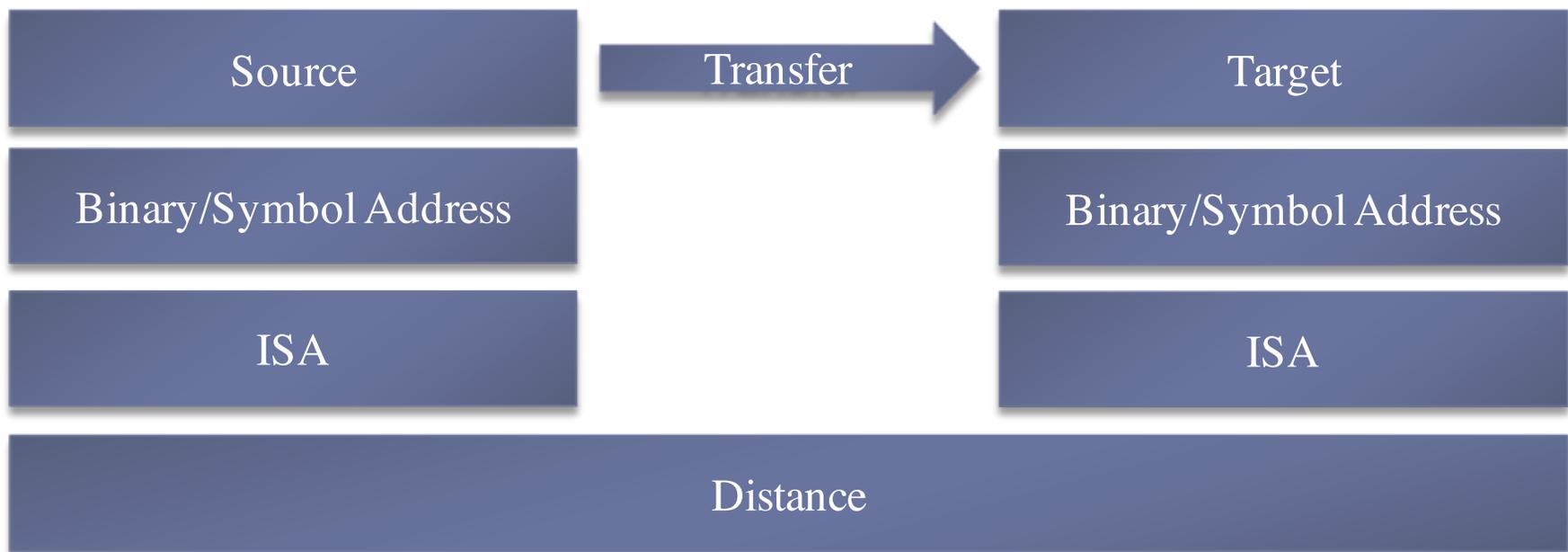
# Branch Transfer Access Control (BTAC)

- \* Inter-components Branch Transfer Access Control
  - \* Branch access per Component
- \* Intra-components Branch Transfer Access Control
  - \* Branch access per branch transfer properties



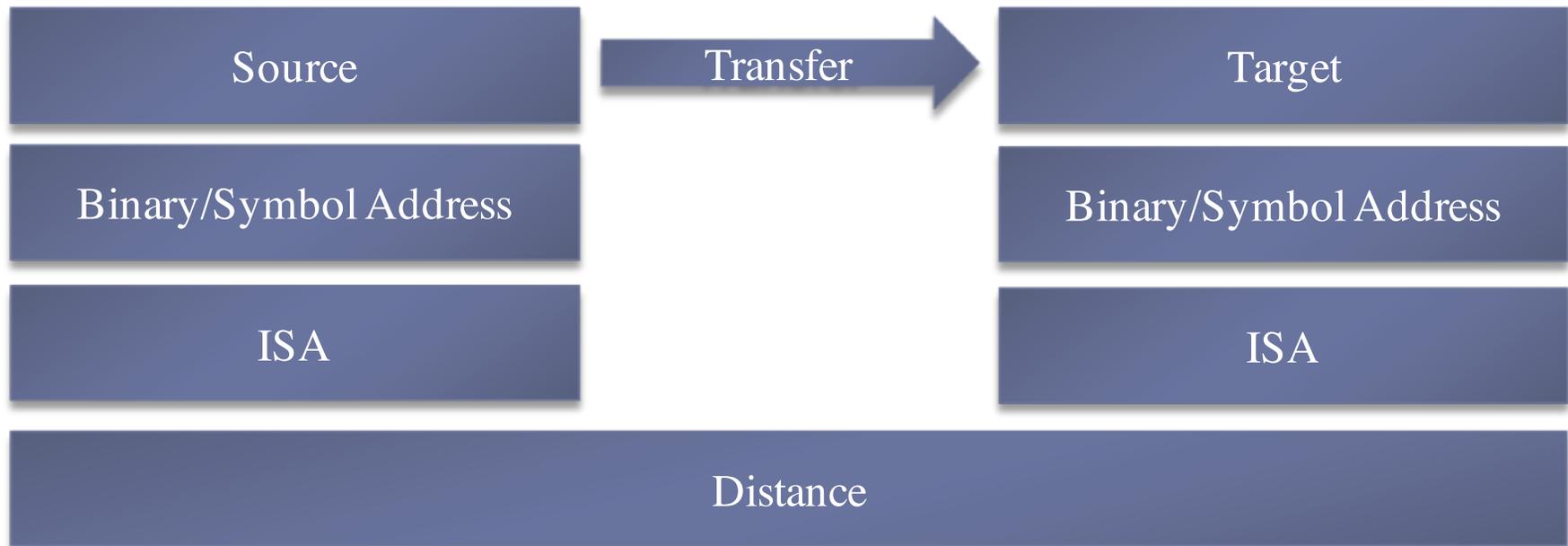
# BTAC vs CFI

- \* BTAC能够使用更多的上下文信息来判定跳转是否合法
- \* CFI仅考虑control flow integrity



# Valid Gadgets

- \* 细粒度CFI再细，也无法阻止valid gadgets
- \* Ret to Call-site, call2libc



# Syscall Access Control (SCAC)

- \* BTAC的特例： Syscall Access Control
  - \* Explicit legal syscall instructions
    - \* Guaranteed by BTAC or CFI
  - \* Access control based on In-process context
    - \* e.g. memory sealing
    - \* More e.g. User-driven Access Control (S&P'12)
  - \* Performance
    - \* Syscall is much less frequent invoked than branch transfer
    - \* Thus we can check more context

# Branch Transfer Access Control with BTF

- \* BTF is the flag in MSR\_DEBUGCTLA MSR
- \* Used to enable single-step on branches

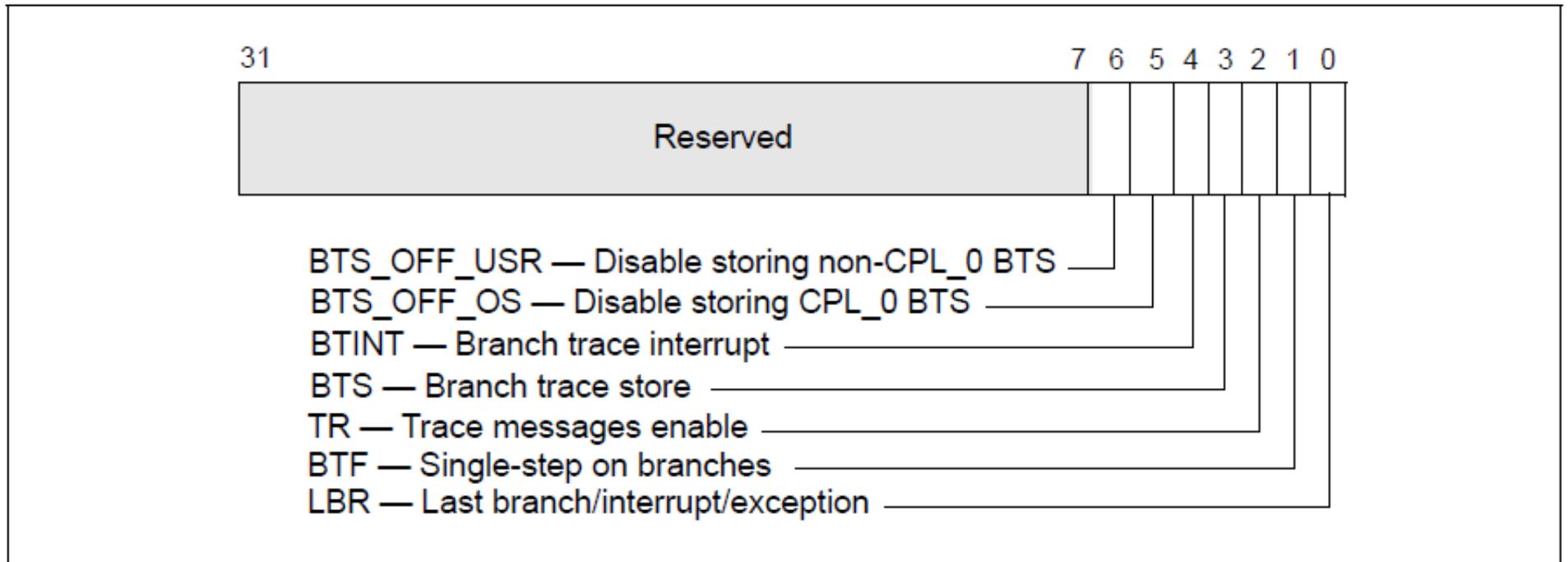
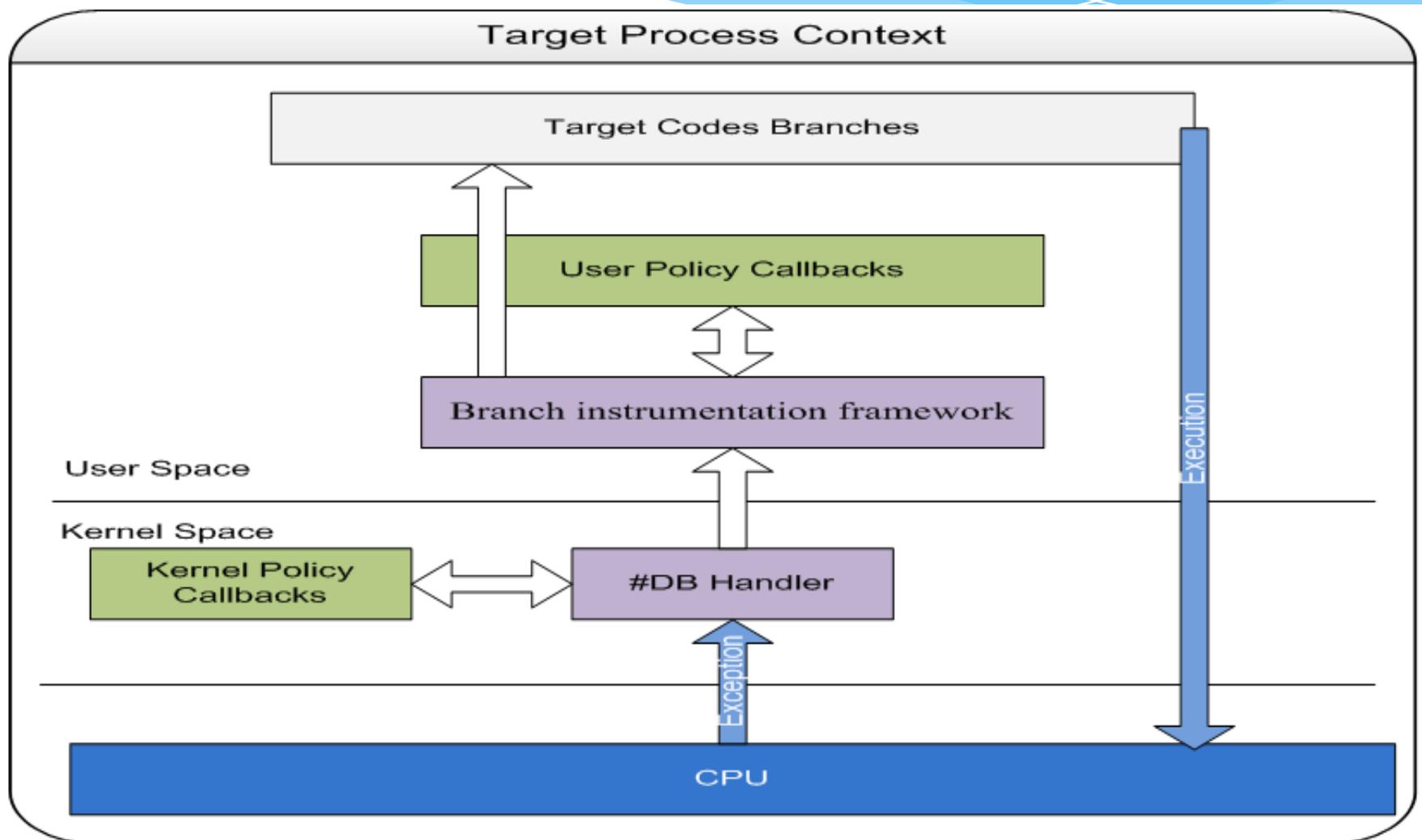


Figure 17-12. MSR\_DEBUGCTLA MSR for Pentium 4 and Intel Xeon Processors

# BTF based BTAC 示例



# BTAC vs. GIFT

- \* GIFT by TK
- \* Internal function should never be called by arbitrary component, for example MSHTML.DLL -> NTDLL
- \* The branch transfers happen between many external components' function and NTDLL internal functions are invalid cases

# BTAC vs. Vital Point Strike

- \* 点穴 by TK
- \* Use Script engine safe mode to call privileged APIs
- \* The branch transfer between script engine and privileged APIs with sensitive contents should be blocked

# BTAC vs. Interdimensional Execution

- \* DVE by Yuange
- \* Use script engine to create arbitrary API calls
- \* JIT code should not directly call into many sensitive APIs

# 结论

- \* CFI is another important security layer, not the final solution
- \* CFI is necessary and effective to support in-process access control, like BTAC, to build a full Airport Security Model
- \* Branch Transfer Access Control is missed part in today's access control facility
- \* BTAC provides solid protections against system level intrusions
  - \* Con: still needs more memory safety support to guarantee in-process abusing
  - \* Pro: don't require full memory safety

# Thanks!



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